Lecture 16 Age-optimal Scheduling in Queues. (3)

Sample-path methods; [] Interchange arguments: W packet 1 2 3 5 Ð Total delay: 3+5=8 H 2 1 2 5 b Total delay: 2+5=7 The total delay reduces by interchanged the two packets. stachastic task arrivals. 2 1 1 1 2 1 Ð 3  $X_1$ : service time of packet 1 = 4 seconds. -2=2 seconds.  $X_2$ : ~

idea: compare the remaining processing time of the packets. Optimal policy for minimizing total delay / average delay: preemptive SRPT (shortest remaining processing time first) 2 forward induction: Define system state.  $\overrightarrow{V_{\pi}(t)}$ , and partial order  $\leq$ goal: prove that, on every Sample peth.  $\overrightarrow{V_{p}}(t) \leq \overrightarrow{V_{\pi}}(t), \quad \forall t, \forall \pi \in \mathcal{T},$ Reading: SRPT 1968. 1978, HW3: Problem 1. Briefly explain the forward indunction method for proving the optimality of preemptive SRPT. Explain, in the proof: () what is the system state? E what is the partial order ?

server scheduler monitor source Server generation time. arrival time, delivery time Packet i Si Cì Di  $0 \in S_1 \leq S_2 \leq -$  $s_i \in C_i \in P_i$ Si and Ci are arbitrarily given. In-order arrivals: SiE Sit, CiE Citi. Uut-of-order arrivals: Si E Si+1. Ci > Ci+1, B: buffer size. M: No. of servers. If B=0, system can keep M packets, i.i.d. non-expenential service times, A01:  $\Delta(t) = t - \max \{ S_i : D_i \leq t \}$ 

New - Better - than - Used (NBU) distribution; X: the service time of a packet. R=[X-t | x>t]: the remaining processing time of a packet after t seconds, given that X > tX Zet R  $Pr(X \ge s) \ge Pr(R \ge s)$  $= \Pr(X-t > s | X > t),$  $= \Pr(X > s + t | X > t)$ Pr(X>s+t)Pr(x>t) $\Pr(X > s+t) \leq \Pr(X > s) \Pr(X > t) \quad \forall s.t \geq 0,$ Examples of NBU distributions: () . exp. distribution. 2 const distribution. 3, const + exp (4). Geo distribution. (3) gamma distribution. Erlang

SOULCE Si quene Ci server Vi destination -D; Pef: Vi: service starting time of packet i.  $S_i \leq C_i \leq V_i \leq D_i$ . Note: For a non-preemptive policy. Di - Vi = service time of packet i. Def: Age of served information W(t),  $W(t) = t - \max\{S_i: \forall i \leq t\}.$  $\Delta(t) = t - \max\{S_i: P_i \leq t\}$  $W(t) \leq \Delta(t)$ 

厶(t) W(t) -1 Si Vi Si+1 D; Viti Piti Ľ Thm: If (i) the packet service times are j.i.d. NBU. then for all  $M \ge 1$ .  $B \ge 1$ , I, and  $\pi \in \pi_{np}$ . [{ Wnon-prmp-LGFt (t) t≥0}] ≤st [{ △n(t), t≥03 ]] where The is the set of non-preemptive policies. or for all non-decreasing functional f.  $E[f(W_{non-prmp-LGFt}(t), t \ge 0]]$  $\leq \min_{\pi \in \pi_{np}} \mathbb{E} \left[ f\left( \left\{ \Delta_{\pi}(t), t \geq 0 \right\} \right) \middle| I \right]$  $E\left[f\left(\left|\Delta_{non-prmp-LGFt}(t), t \ge 0\right>\right| \Gamma\right]$ Ś

 $E.g.: f(\{\Delta(t), t \ge 0\}) = \limsup_{T \to \infty} \frac{1}{T} \int_{0}^{t} \Delta(t) dt,$ Corollary: Under the conditions of the Thm, we have  $\begin{array}{c|c} \min & \limsup & \frac{1}{T} \in \left[ \int_{0}^{T} \Delta_{\pi}(t) dt \right] \\ \overline{\Pi} \in \overline{\Pi}_{np} & T \rightarrow \infty \end{array}$ < limsup - E [ [ Anon-prmp-LGFS (+) dt ]] T> 00  $\leq \min_{\pi \in \pi_{np}} \lim_{T \to \infty} \frac{1}{T} E\left[\int_{0}^{T} \Delta_{\pi}(t) dt | I\right] + E[X]$ mean service time Proof of Thm: Section 2.6.2 of the book Sun, Kadota, Talak Modiano 2019, HW3. Problem 2. Read Section 2.6.2 explain your understanding of weak work-efficiency ordering" and its role in the proof of the Thm.